

- An attempt to access memory with an effective address alignment that is invalid for the instruction causes the alignment exception handler to be invoked.
- The execution of an **sc** instruction permits a program to call on the system to perform a service, by causing a system call exception handler to be invoked.
- The execution of a trap instruction invokes the program exception trap handler.
- The execution of a floating-point instruction when floating-point instructions are disabled invokes the floating-point unavailable exception handler.
- The execution of an instruction that causes a floating-point exception that is enabled invokes the floating-point enabled exception handler.
- The execution of a floating-point instruction that requires system software assistance causes the floating-point assist exception handler to be invoked. The conditions under which such software assistance is required are implementation-dependent.

Exceptions caused by asynchronous events are described in Chapter 6, “Exceptions.”

4.2 PowerPC UISA Instructions

The PowerPC user instruction set architecture (UISA) includes the base user-level instruction set (excluding a few user-level cache-control, synchronization, and time base instructions), user-level registers, programming model, data types, and addressing modes. This section discusses the instructions defined in the UISA.

4.2.1 Integer Instructions

The integer instructions consist of the following:

- Integer arithmetic instructions
- Integer compare instructions
- Integer logical instructions
- Integer rotate and shift instructions

Integer instructions use the content of the GPRs as source operands and place results into GPRs. Integer arithmetic, shift, rotate, and string move instructions may update or read values from the XER, and the condition register (CR) fields may be updated if the Rc bit of the instruction is set.

These instructions treat the source operands as signed integers unless the instruction is explicitly identified as performing an unsigned operation. For example, Multiply High-Word Unsigned (**mulhwu**) and Divide Word Unsigned (**divwu**) instructions interpret both operands as unsigned integers.

The integer instructions that are coded to update the condition register, and the integer arithmetic instruction, **addic.**, set CR bits 0–3 (CR0) to characterize the result of the operation. CR0 is set to reflect a signed comparison of the result to zero.

The integer arithmetic instructions, **addic**, **addic.**, **subfic**, **addc**, **subfc**, **adde**, **subfe**, **addme**, **subme**, **addze**, and **subfze**, always set the XER bit, CA, to reflect the carry out of bit 0. Integer arithmetic instructions with the overflow enable (OE) bit set in the instruction encoding (instructions with o suffix) cause the XER[SO] and XER[OV] to reflect an overflow of the result. Except for the multiply low and divide instructions, these integer arithmetic instructions reflect the overflow of the result.

Instructions that select the overflow option (enable XER[OV]) or that set the XER carry bit (CA) may delay the execution of subsequent instructions.

Unless otherwise noted, when CR0 and the XER are set, they reflect the value placed in the target register.

4.2.1.1 Integer Arithmetic Instructions

Table 4-1 lists the integer arithmetic instructions for the PowerPC processors.

Table 4-1. Integer Arithmetic Instructions

Name	Mnemonic	Operand Syntax	Operation
Add Immediate	addi	rD,rA,SIMM	The sum (rA 0) + SIMM is placed into rD.
Add Immediate Shifted	addis	rD,rA,SIMM	The sum (rA 0) + (SIMM 0x0000) is placed into rD.
Add	add add. addo addo.	rD,rA,rB	The sum (rA) + (rB) is placed into rD. add Add add. Add with CR Update. The dot suffix enables the update of the CR. addo Add with Overflow Enabled. The o suffix enables the overflow bit (OV) in the XER. addo. Add with Overflow and CR Update. The o. suffix enables the update of the CR and enables the overflow bit (OV) in the XER.
Subtract From	subf subf. subfo subfo.	rD,rA,rB	The sum \neg (rA) + (rB) + 1 is placed into rD. subf Subtract From subf. Subtract from with CR Update. The dot suffix enables the update of the CR. subfo Subtract from with Overflow Enabled. The o suffix enables the overflow bit (OV) in the XER. subfo. Subtract from with Overflow and CR Update. The o. suffix enables the update of the CR and enables the overflow bit (OV) in the XER.
Add Immediate Carrying	addic	rD,rA,SIMM	The sum (rA) + SIMM is placed into rD.
Add Immediate Carrying and Record	addic.	rD,rA,SIMM	The sum (rA) + SIMM is placed into rD. The CR is updated.
Subtract from Immediate Carrying	subfic	rD,rA,SIMM	The sum \neg (rA) + SIMM + 1 is placed into rD.

Table 4-1. Integer Arithmetic Instructions (Continued)

Name	Mnemonic	Operand Syntax	Operation
Add Carrying	addc addc. addco addco.	rD,rA,rB	The sum $(rA) + (rB)$ is placed into rD. addc Add Carrying addc. Add Carrying with CR Update. The dot suffix enables the update of the CR. addco Add Carrying with Overflow Enabled. The o suffix enables the overflow bit (OV) in the XER. addco. Add Carrying with Overflow and CR Update. The o. suffix enables the update of the CR and enables the overflow bit (OV) in the XER.
Subtract from Carrying	subfc subfc. subfco subfco.	rD,rA,rB	The sum $\neg(rA) + (rB) + 1$ is placed into rD. subfc Subtract from Carrying subfc. Subtract from Carrying with CR Update. The dot suffix enables the update of the CR. subfco Subtract from Carrying with Overflow. The o suffix enables the overflow bit (OV) in the XER. subfco. Subtract from Carrying with Overflow and CR Update. The o. suffix enables the update of the CR and enables the overflow bit (OV) in the XER.
Add Extended	adde adde. addeo addeo.	rD,rA,rB	The sum $(rA) + (rB) + XER[CA]$ is placed into rD. adde Add Extended adde. Add Extended with CR Update. The dot suffix enables the update of the CR. addeo Add Extended with Overflow. The o suffix enables the overflow bit (OV) in the XER. addeo. Add Extended with Overflow and CR Update. The o. suffix enables the update of the CR and enables the overflow bit (OV) in the XER.
Subtract from Extended	subfe subfe. subfeo subfeo.	rD,rA,rB	The sum $\neg(rA) + (rB) + XER[CA]$ is placed into rD. subfe Subtract from Extended subfe. Subtract from Extended with CR Update. The dot suffix enables the update of the CR. subfeo Subtract from Extended with Overflow. The o suffix enables the overflow bit (OV) in the XER. subfeo. Subtract from Extended with Overflow and CR Update. The o. suffix enables the update of the CR and enables the overflow (OV) bit in the XER.
Add to Minus One Extended	addme addme. addmeo addmeo.	rD,rA	The sum $(rA) + XER[CA]$ added to 0xFFFF_FFFF is placed into rD. addme Add to Minus One Extended addme. Add to Minus One Extended with CR Update. The dot suffix enables the update of the CR. addmeo Add to Minus One Extended with Overflow. The o suffix enables the overflow bit (OV) in the XER. addmeo. Add to Minus One Extended with Overflow and CR Update. The o. suffix enables the update of the CR and enables the overflow (OV) bit in the XER.

Table 4-1. Integer Arithmetic Instructions (Continued)

Name	Mnemonic	Operand Syntax	Operation
Subtract from Minus One Extended	subfme subfme. subfmeo subfmeo.	rD,rA	The sum $\neg (rA) + XER[CA]$ added to 0xFFFF_FFFF is placed into rD. subfme Subtract from Minus One Extended subfme. Subtract from Minus One Extended with CR Update. The dot suffix enables the update of the CR. subfmeo Subtract from Minus One Extended with Overflow. The o suffix enables the overflow bit (OV) in the XER. subfmeo. Subtract from Minus One Extended with Overflow and CR Update. The o. suffix enables the update of the CR and enables the overflow bit (OV) in the XER.
Add to Zero Extended	addze addze. addzeo addzeo.	rD,rA	The sum $(rA) + XER[CA]$ is placed into rD. addze Add to Zero Extended addze. Add to Zero Extended with CR Update. The dot suffix enables the update of the CR. addzeo Add to Zero Extended with Overflow. The o suffix enables the overflow bit (OV) in the XER. addzeo. Add to Zero Extended with Overflow and CR Update. The o. suffix enables the update of the CR and enables the overflow bit (OV) in the XER.
Subtract from Zero Extended	subfze subfze. subfzeo subfzeo.	rD,rA	The sum $\neg (rA) + XER[CA]$ is placed into rD. subfze Subtract from Zero Extended subfze. Subtract from Zero Extended with CR Update. The dot suffix enables the update of the CR. subfzeo Subtract from Zero Extended with Overflow. The o suffix enables the overflow bit (OV) in the XER. subfzeo. Subtract from Zero Extended with Overflow and CR Update. The o. suffix enables the update of the CR and enables the overflow bit (OV) in the XER.
Negate	neg neg. nego nego.	rD,rA	The sum $\neg (rA) + 1$ is placed into rD. neg Negate neg. Negate with CR Update. The dot suffix enables the update of the CR. nego Negate with Overflow. The o suffix enables the overflow bit (OV) in the XER. nego. Negate with Overflow and CR Update. The o. suffix enables the update of the CR and enables the overflow bit (OV) in the XER.
Multiply Low Immediate	mulli	rD,rA,SIMM	The low-order 32 bits of the product $(rA) * SIMM$ are placed into rD. This instruction can be used with mulhdx or mulhwX to calculate a full 64-bit product.

Table 4-1. Integer Arithmetic Instructions (Continued)

Name	Mnemonic	Operand Syntax	Operation
Multiply Low	mullw mullw. mullwo mullwo.	rD,rA,rB	<p>The 32-bit product (rA) * (rB) is placed into register rD.</p> <p>This instruction can be used with mulhwx to calculate a full 64-bit product.</p> <p>mullw Multiply Low</p> <p>mullw. Multiply Low with CR Update. The dot suffix enables the update of the CR.</p> <p>mullwo Multiply Low with Overflow. The o suffix enables the overflow bit (OV) in the XER.</p> <p>mullwo. Multiply Low with Overflow and CR Update. The o. suffix enables the update of the condition register and enables the overflow bit (OV) in the XER.</p>
Multiply High Word	mulhw mulhw.	rD,rA,rB	<p>The contents of rA and rB are interpreted as 32-bit signed integers. The 64-bit product is formed. The high-order 32 bits of the 64-bit product are placed into rD.</p> <p>mulhw Multiply High Word</p> <p>mulhw. Multiply High Word with CR Update. The dot suffix enables the update of the CR.</p>
Multiply High Word Unsigned	mulhwu mulhwu.	rD,rA,rB	<p>The contents of rA and of rB are interpreted as 32-bit unsigned integers. The 64-bit product is formed. The high-order 32 bits of the 64-bit product are placed into rD.</p> <p>mulhwu Multiply High Word Unsigned</p> <p>mulhwu. Multiply High Word Unsigned with CR Update. The dot suffix enables the update of the CR.</p>
Divide Word	divw divw. divwo divwo.	rD,rA,rB	<p>The dividend is the signed value of rA. The divisor is the signed value of rB. The quotient is placed into rD. The remainder is not supplied as a result.</p> <p>divw Divide Word</p> <p>divw. Divide Word with CR Update. The dot suffix enables the update of the CR.</p> <p>divwo Divide Word with Overflow. The o suffix enables the overflow bit (OV) in the XER.</p> <p>divwo. Divide Word with Overflow and CR Update. The o. suffix enables the update of the CR and enables the overflow bit (OV) in the XER.</p>
Divide Word Unsigned	divwu divwu. divwuo divwuo.	rD,rA,rB	<p>The dividend is the zero-extended value in rA. The divisor is the zero-extended value in rB. The quotient is placed into rD. The remainder is not supplied as a result.</p> <p>divwu Divide Word Unsigned</p> <p>divwu. Divide Word Unsigned with CR Update. The dot suffix enables the update of the CR.</p> <p>divwuo Divide Word Unsigned with Overflow. The o suffix enables the overflow bit (OV) in the XER.</p> <p>divwuo. Divide Word Unsigned with Overflow and CR Update. The o. suffix enables the update of the CR and enables the overflow bit (OV) in the XER.</p>

Although there is no “Subtract Immediate” instruction, its effect can be achieved by using an **addi** instruction with the immediate operand negated. Simplified mnemonics are provided that include this negation. The **subf** instructions subtract the second operand (**rA**) from the third operand (**rB**). Simplified mnemonics are provided in which the third operand is subtracted from the second operand. See Appendix F, “Simplified Mnemonics,” for examples.

4.2.1.2 Integer Compare Instructions

The integer compare instructions algebraically or logically compare the contents of register **rA** with either the zero-extended value of the UIMM operand, the sign-extended value of the SIMM operand, or the contents of register **rB**. The comparison is signed for the **cmpi** and **cmp** instructions, and unsigned for the **cmpli** and **cmpl** instructions. Table 4-2 summarizes the integer compare instructions.

Appendix F, “Simplified Mnemonics” For 32-bit implementations, the L field must be cleared, otherwise the instruction form is invalid.

The integer compare instructions (shown in Table 4-2) set one of the leftmost three bits of the designated CR field, and clear the other two. XER[SO] is copied into bit 3 of the CR field.

Table 4-2. Integer Compare Instructions

Name	Mnemonic	Operand Syntax	Operation
Compare Immediate	cmpi	crfD ,L,rA,SIMM	The value in register rA is compared with the sign-extended value of the SIMM operand, treating the operands as signed integers. The result of the comparison is placed into the CR field specified by operand crfD .
Compare	cmp	crfD ,L,rA,rB	The value in register rA is compared with the value in register rB , treating the operands as signed integers. The result of the comparison is placed into the CR field specified by operand crfD .
Compare Logical Immediate	cmpli	crfD ,L,rA,UIMM	The value in register rA is compared with 0x0000 UIMM, treating the operands as unsigned integers. The result of the comparison is placed into the CR field specified by operand crfD .
Compare Logical	cmpl	crfD ,L,rA,rB	The value in register rA is compared with the value in register rB , treating the operands as unsigned integers. The result of the comparison is placed into the CR field specified by operand crfD .

The **crfD** operand can be omitted if the result of the comparison is to be placed in CR0. Otherwise the target CR field must be specified in the instruction **crfD** field, using an explicit field number.

For information on simplified mnemonics for the integer compare instructions see Appendix F, “Simplified Mnemonics.”

4.2.1.3 Integer Logical Instructions

The logical instructions shown in Table 4-3 perform bit-parallel operations on 32-bit operands. Logical instructions with the CR updating enabled (uses dot suffix) and instructions **andi.** and **andis.** set CR field CR0 (bits 0 to 2) to characterize the result of the logical operation. Logical instructions without CR update and the remaining logical instructions do not modify the CR. Logical instructions do not affect the XER[SO], XER[OV], and XER[CA] bits.

See Appendix F, “Simplified Mnemonics,” for simplified mnemonic examples for integer logical operations.

Table 4-3. Integer Logical Instructions

Name	Mnemonic	Operand Syntax	Operation
AND Immediate	andi.	rA,rS,UIMM	The contents of rS are ANDed with 0x0000 UIMM and the result is placed into rA. The CR is updated.
AND Immediate Shifted	andis.	rA,rS,UIMM	The content of rS are ANDed with UIMM 0x0000 and the result is placed into rA. The CR is updated.
OR Immediate	ori	rA,rS,UIMM	The contents of rS are ORed with 0x0000 UIMM and the result is placed into rA. The preferred no-op is ori 0,0,0
OR Immediate Shifted	oris	rA,rS,UIMM	The contents of rS are ORed with UIMM 0x0000 and the result is placed into rA.
XOR Immediate	xori	rA,rS,UIMM	The contents of rS are XORed with 0x0000 UIMM and the result is placed into rA.
XOR Immediate Shifted	xoris	rA,rS,UIMM	The contents of rS are XORed with UIMM 0x0000 and the result is placed into rA.
AND	and and.	rA,rS,rB	The contents of rS are ANDed with the contents of register rB and the result is placed into rA. and AND and. AND with CR Update. The dot suffix enables the update of the CR.
OR	or or.	rA,rS,rB	The contents of rS are ORed with the contents of rB and the result is placed into rA. or OR or. OR with CR Update. The dot suffix enables the update of the CR.
XOR	xor xor.	rA,rS,rB	The contents of rS are XORed with the contents of rB and the result is placed into rA. xor XOR xor. XOR with CR Update. The dot suffix enables the update of the CR.

Table 4-3. Integer Logical Instructions (Continued)

Name	Mnemonic	Operand Syntax	Operation
NAND	nand nand.	rA,rS,rB	The contents of rS are ANDed with the contents of rB and the one's complement of the result is placed into rA. nand NAND nand. NAND with CR Update. The dot suffix enables the update of CR. Note that nandx , with rS = rB, can be used to obtain the one's complement.
NOR	nor nor.	rA,rS,rB	The contents of rS are ORed with the contents of rB and the one's complement of the result is placed into rA. nor NOR nor. NOR with CR Update. The dot suffix enables the update of the CR. Note that norx , with rS = rB, can be used to obtain the one's complement.
Equivalent	eqv eqv.	rA,rS,rB	The contents of rS are XORed with the contents of rB and the complemented result is placed into rA. eqv Equivalent eqv. Equivalent with CR Update. The dot suffix enables the update of the CR.
AND with Complement	andc andc.	rA,rS,rB	The contents of rS are ANDed with the one's complement of the contents of rB and the result is placed into rA. andc AND with Complement andc. AND with Complement with CR Update. The dot suffix enables the update of the CR.
OR with Complement	orc orc.	rA,rS,rB	The contents of rS are ORed with the complement of the contents of rB and the result is placed into rA. orc OR with Complement orc. OR with Complement with CR Update. The dot suffix enables the update of the CR.
Extend Sign Byte	extsb extsb.	rA,rS	The contents of the low-order eight bits of rS are placed into the low-order eight bits of rA. Bit 24 of rS is placed into the remaining high-order bits of rA. extsb Extend Sign Byte extsb. Extend Sign Byte with CR Update. The dot suffix enables the update of the CR.
Extend Sign Half Word	extsh extsh.	rA,rS	The contents of the low-order 16 bits of rS are placed into the low-order 16 bits of rA. Bit 16 of rS is placed into the remaining high-order bits of rA. extsh Extend Sign Half Word extsh. Extend Sign Half Word with CR Update. The dot suffix enables the update of the CR.
Count Leading Zeros Word	cntlzw cntlzw.	rA,rS	A count of the number of consecutive zero bits starting at bit 0 of rS is placed into rA. This number ranges from 0 to 32, inclusive. If Rc = 1 (dot suffix), LT is cleared in CR0. cntlzw Count Leading Zeros Word cntlzw. Count Leading Zeros Word with CR Update. The dot suffix enables the update of the CR.

4.2.1.4 Integer Rotate and Shift Instructions

Rotation operations are performed on data from a GPR, and the result, or a portion of the result, is returned to a GPR. The rotation operations rotate a 32-bit quantity left by a specified number of bit positions. Bits that exit from position 0 enter at position 31.

The rotate and shift instructions employ a mask generator. The mask is 32 bits long and consists of '1' bits from a start bit, *Mstart*, through and including a stop bit, *Mstop*, and '0' bits elsewhere. The values of *Mstart* and *Mstop* range from 0 to 31. If *Mstart* > *Mstop*, the '1' bits wrap around from position 31 to position 0. Thus the mask is formed as follows:

if $Mstart \leq Mstop$ then

 mask[mstart–mstop] = ones

 mask[all other bits] = zeros

else

 mask[mstart–31] = ones

 mask[0–mstop] = ones

 mask[all other bits] = zeros

It is not possible to specify an all-zero mask. The use of the mask is described in the following sections.

If CR updating is enabled, rotate and shift instructions set CR0[0–2] according to the contents of *rA* at the completion of the instruction. Rotate and shift instructions do not change the values of XER[OV] and XER[SO] bits. Rotate and shift instructions, except algebraic right shifts, do not change the XER[CA] bit.

See Appendix F, “Simplified Mnemonics,” for a complete list of simplified mnemonics that allows simpler coding of often-used functions such as clearing the leftmost or rightmost bits of a register, left justifying or right justifying an arbitrary field, and simple rotates and shifts.

4.2.1.4.1 Integer Rotate Instructions

Integer rotate instructions rotate the contents of a register. The result of the rotation is either inserted into the target register under control of a mask (if a mask bit is 1 the associated bit of the rotated data is placed into the target register, and if the mask bit is 0 the associated bit in the target register is unchanged), or ANDed with a mask before being placed into the target register.

Rotate left instructions allow right-rotation of the contents of a register to be performed by a left-rotation of $64 - n$, where n is the number of bits by which to rotate right. It also allows right-rotation of the contents of the low-order 32 bits of a register to be performed by a left-rotation of $32 - n$, where n is the number of bits by which to rotate right.

The integer rotate instructions are summarized in Table 4-4.

Table 4-4. Integer Rotate Instructions

Name	Mnemonic	Operand Syntax	Operation
Rotate Left Word Immediate then AND with Mask	rlwinm rlwinm.	rA,rS,SH,MB,ME	The contents of register rS are rotated left by the number of bits specified by operand SH. A mask is generated having 1 bits from the bit specified by operand MB through the bit specified by operand ME and 0 bits elsewhere. The rotated data is ANDed with the generated mask and the result is placed into register rA. rlwinm Rotate Left Word Immediate then AND with Mask rlwinm. Rotate Left Word Immediate then AND with Mask with CR Update. The dot suffix enables the update of the CR.
Rotate Left Word then AND with Mask	rlwnm rlwnm.	rA,rS,rB,MB,ME	The contents of rS are rotated left by the number of bits specified by operand in the low-order five bits of rB. A mask is generated having 1 bits from the bit specified by operand MB through the bit specified by operand ME and 0 bits elsewhere. The rotated word is ANDed with the generated mask and the result is placed into rA. rlwnm Rotate Left Word then AND with Mask rlwnm. Rotate Left Word then AND with Mask with CR Update. The dot suffix enables the update of the CR.
Rotate Left Word Immediate then Mask Insert	rlwimi rlwimi.	rA,rS,SH,MB,ME	The contents of rS are rotated left by the number of bits specified by operand SH. A mask is generated having 1 bits from the bit specified by operand MB through the bit specified by operand ME and 0 bits elsewhere. The rotated word is inserted into rA under control of the generated mask. rlwimi Rotate Left Word Immediate then Mask rlwimi. Rotate Left Word Immediate then Mask Insert with CR Update. The dot suffix enables the update of the CR.

4.2.1.4.2 Integer Shift Instructions

The integer shift instructions perform left and right shifts. Immediate-form logical (unsigned) shift operations are obtained by specifying masks and shift values for certain rotate instructions. Simplified mnemonics (shown in Appendix F, “Simplified Mnemonics”) are provided to make coding of such shifts simpler and easier to understand.

Any shift right algebraic instruction, followed by **addze**, can be used to divide quickly by 2^n . The setting of XER[CA] by the shift right algebraic instruction is independent of mode.

Multiple-precision shifts can be programmed as shown in Appendix C, “Multiple-Precision Shifts.”

The integer shift instructions are summarized in Table 4-5.

Table 4-5. Integer Shift Instructions

Name	Mnemonic	Operand Syntax	Operation
Shift Left Word	slw slw.	rA,rS,rB	The contents of rS are shifted left the number of bits specified by operand in the low-order six bits of rB. Bits shifted out of position 0 are lost. Zeros are supplied to the vacated positions on the right. The 32-bit result is placed into rA. slw Shift Left Word slw. Shift Left Word with CR Update. The dot suffix enables the update of the CR.
Shift Right Word	srw srw.	rA,rS,rB	The contents of rS are shifted right the number of bits specified by the low-order six bits of rB. Bits shifted out of position 31 are lost. Zeros are supplied to the vacated positions on the left. The 32-bit result is placed into rA. srw Shift Right Word srw. Shift Right Word with CR Update. The dot suffix enables the update of the CR.
Shift Right Algebraic Word Immediate	srawi srawi.	rA,rS,SH	The contents of rS are shifted right the number of bits specified by operand SH. Bits shifted out of position 31 are lost. The result is sign extended and placed into rA. srawi Shift Right Algebraic Word Immediate srawi. Shift Right Algebraic Word Immediate with CR Update. The dot suffix enables the update of the CR.
Shift Right Algebraic Word	sraw sraw.	rA,rS,rB	The contents of rS are shifted right the number of bits specified by the low-order six bits of rB. Bits shifted out of position 31 are lost. The result is placed into rA. sraw Shift Right Algebraic Word sraw. Shift Right Algebraic Word with CR Update. The dot suffix enables the update of the CR.

4.2.2 Floating-Point Instructions

This section describes the floating-point instructions, which include the following:

- Floating-point arithmetic instructions
- Floating-point multiply-add instructions
- Floating-point rounding and conversion instructions
- Floating-point compare instructions
- Floating-point status and control register instructions
- Floating-point move instructions

Note that MSR[FP] must be set in order for any of these instructions (including the floating-point loads and stores) to be executed. If MSR[FP] = 0 when any floating-point instruction is attempted, the floating-point unavailable exception is taken (see Section 6.4.8, “Floating-Point Unavailable Exception (0x00800)”). See Section 4.2.3, “Load and Store Instructions,” for information about floating-point loads and stores.