## EECS483 D13: SSA Example

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## Announcements

- Homework 5 on CTools
-Due on 4/22


## Static Single Assignment Form

- Each variable is given a unique name when it is assigned to a new value
- All of the uses of this assignment are renamed accordingly
- Phi nodes: a special multiplexer that choose a value from its arguments


## SSA Conversion (1/2)



## SSA Conversion (2/2)

- Dominator analysis
-Find the dominator frontier set DF(BB) for each basic block BB
- Phi node insertion
-If variable $x$ is defined in $B B$, then a Phi node of $x$ is needed in each basic block in DF(BB)
- Variable renaming
-Rename variables in each assignment (including Phi node) and all their uses


## Dominator Analysis (1/2)

- X dominates Y if every path from entry to $Y$ contains $X$
-X dominates X itself
- $Z$ is a dominance frontier of $X$ if $X$ dominates a predecessor $Y$ of $Z$ but not Z
-The first BB that is not dominated by X
- If variable $a$ is defined in $X$
-Uses of a in $Y$ refer to the definition in $X$
-Uses of a in $Z$ don't necessary refer to the definition in $X$

- Need a Phi node for a!


## Dominator Analysis (2/2)

- $\operatorname{Dom}(X)=$ Intersection(Dom(predecessors of $X)$ )
- Compute dominators
- Initialization
- Dom(Entry) = \{Entry $\}$
- $\operatorname{Dom}(X)=\{$ all nodes $\}$ for all other $X$
-While(change):
- Update Dom(X) for each X
- Compute dominance frontiers
-for each Z
- for each predecessor $Y$ of $Z$
- for each $X$ in $\operatorname{Dom}(Y)$ - $\operatorname{Dom}(Z)$
» Put $Z$ into $D F(X)$


## Dominator Analysis: Example



| BB | Dom | DF |
| :---: | :---: | :---: |
| 0 |  |  |
| 1 |  |  |
| 2 |  |  |
| 3 |  |  |
| 4 |  |  |
| 5 |  |  |
| 6 |  |  |
| 7 |  |  |

This example comes from Prof. Mahlke's EECS583 slides.

## Dominator Analysis: Example



| BB | Dom | DF |
| :---: | :--- | :--- |
| 0 | 0 |  |
| 1 | 0,1 |  |
| 2 | $0,1,2$ |  |
| 3 | $0,1,3$ |  |
| 4 | $0,1,3,4$ |  |
| 5 | $0,1,3,5$ |  |
| 6 | $0,1,3,6$ |  |
| 7 | $0,1,7$ |  |

## Dominator Analysis: Example



| BB | Dom | DF |
| :---: | :--- | :--- |
| 0 | 0 | - |
| 1 | 0,1 | - |
| 2 | $0,1,2$ | 7 |
| 3 | $0,1,3$ | 7 |
| 4 | $0,1,3,4$ | 6 |
| 5 | $0,1,3,5$ | 6 |
| 6 | $0,1,3,6$ | 7 |
| 7 | $0,1,7$ | 1 |

## Phi Node Insertion

- Liveness analysis
$-\operatorname{IN}(\mathrm{BB})$ : variables used in BB but defined elsewhere
$-K I L L(B B)$ : variables defined in $B B$
- Algorithm
-for each variable $v$ in $\operatorname{IN}(B B)$ for some $B B$
- $\operatorname{Def}(v)=\{B B: v \in \operatorname{KILL}(B B)\}$
- for each BB $\in$ in $\operatorname{Def}(v)$
- Insert a Phi node for a in DF(BB)
- Add BB into Def(v)


## Phi Node Insertion: Example



| BB | Dom | DF |
| :---: | :--- | :--- |
| 0 | 0 | - |
| 1 | 0,1 | - |
| 2 | $0,1,2$ | 7 |
| 3 | $0,1,3$ | 7 |
| 4 | $0,1,3,4$ | 6 |
| 5 | $0,1,3,5$ | 6 |
| 6 | $0,1,3,6$ | 7 |
| 7 | $0,1,7$ | 1 |


| variable | Def |
| :---: | :---: |
| $a$ |  |
| $b$ |  |
| c |  |
| $d$ |  |
| $i$ |  |

## Phi Node Insertion: Example



| BB | Dom | DF |
| :---: | :--- | :--- |
| 0 | 0 | - |
| 1 | 0,1 | - |
| 2 | $0,1,2$ | 7 |
| 3 | $0,1,3$ | 7 |
| 4 | $0,1,3,4$ | 6 |
| 5 | $0,1,3,5$ | 6 |
| 6 | $0,1,3,6$ | 7 |
| 7 | $0,1,7$ | 1 |


| variable | Def |
| :---: | :---: |
| $a$ | $0,1,3$ |
| $b$ |  |
| $c$ |  |
| $d$ |  |
| $i$ |  |

## Phi Node Insertion: Example



| BB | Dom | DF |
| :---: | :--- | :--- |
| 0 | 0 | - |
| 1 | 0,1 | - |
| 2 | $0,1,2$ | 7 |
| 3 | $0,1,3$ | 7 |
| 4 | $0,1,3,4$ | 6 |
| 5 | $0,1,3,5$ | 6 |
| 6 | $0,1,3,6$ | 7 |
| 7 | $0,1,7$ | 1 |


| variable | Def |
| :---: | :---: |
| $a$ | $0,1,3,7$ |
| $b$ |  |
| $c$ |  |
| $d$ |  |
| $i$ |  |

## Phi Node Insertion: Example



| BB | Dom | DF |
| :---: | :--- | :--- |
| 0 | 0 | - |
| 1 | 0,1 | - |
| 2 | $0,1,2$ | 7 |
| 3 | $0,1,3$ | 7 |
| 4 | $0,1,3,4$ | 6 |
| 5 | $0,1,3,5$ | 6 |
| 6 | $0,1,3,6$ | 7 |
| 7 | $0,1,7$ | 1 |


| variable | Def |
| :---: | :---: |
| $a$ | $0,1,3,7$ |
| $b$ |  |
| $c$ |  |
| $d$ |  |
| $i$ |  |

## Phi Node Insertion: Example



| BB | Dom | DF |
| :---: | :--- | :--- |
| 0 | 0 | - |
| 1 | 0,1 | - |
| 2 | $0,1,2$ | 7 |
| 3 | $0,1,3$ | 7 |
| 4 | $0,1,3,4$ | 6 |
| 5 | $0,1,3,5$ | 6 |
| 6 | $0,1,3,6$ | 7 |
| 7 | $0,1,7$ | 1 |


| variable | Def |
| :---: | :--- |
| $a$ | $0,1,3,7$ |
| $b$ | $0,2,6$ |
| $c$ |  |
| $d$ |  |
| $i$ |  |

## Phi Node Insertion: Example



| BB | Dom | DF |
| :---: | :--- | :--- |
| 0 | 0 | - |
| 1 | 0,1 | - |
| 2 | $0,1,2$ | 7 |
| 3 | $0,1,3$ | 7 |
| 4 | $0,1,3,4$ | 6 |
| 5 | $0,1,3,5$ | 6 |
| 6 | $0,1,3,6$ | 7 |
| 7 | $0,1,7$ | 1 |


| variable | Def |
| :---: | :--- |
| $a$ | $0,1,3,7$ |
| $b$ | $0,2,6,7$ |
| $c$ |  |
| $d$ |  |
| $i$ |  |
|  |  |

## Phi Node Insertion: Example



| BB | Dom | DF |
| :---: | :--- | :--- |
| 0 | 0 | - |
| 1 | 0,1 | - |
| 2 | $0,1,2$ | 7 |
| 3 | $0,1,3$ | 7 |
| 4 | $0,1,3,4$ | 6 |
| 5 | $0,1,3,5$ | 6 |
| 6 | $0,1,3,6$ | 7 |
| 7 | $0,1,7$ | 1 |


| variable | Def |
| :---: | :--- |
| $a$ | $0,1,3,7$ |
| $b$ | $0,2,6,7,1$ |
| $c$ |  |
| $d$ |  |
| $i$ |  |

## Phi Node Insertion: Example



| BB | Dom | DF |
| :---: | :--- | :--- |
| 0 | 0 | - |
| 1 | 0,1 | - |
| 2 | $0,1,2$ | 7 |
| 3 | $0,1,3$ | 7 |
| 4 | $0,1,3,4$ | 6 |
| 5 | $0,1,3,5$ | 6 |
| 6 | $0,1,3,6$ | 7 |
| 7 | $0,1,7$ | 1 |


| variable | Def |
| :---: | :--- |
| $a$ | $0,1,3,7$ |
| $b$ | $0,2,6,7,1$ |
| $c$ | $0,1,2,5$ |
| $d$ |  |
| $i$ |  |

## Phi Node Insertion: Example



## Phi Node Insertion: Example



## Phi Node Insertion: Example



## Phi Node Insertion: Example



## Phi Node Insertion: Example



## Phi Node Insertion: Example



## Phi Node Insertion: Example



## Variable Renaming (1/3)

- Constructing the dominator tree
-The parent of a basic block is its immediate dominator
- For each variable, maintain the following data structures
-A counter for creating new names
-A stack to keep track of currently available names for this variable
- The top of the stack is the name defined in its nearest dominators



## Variable Renaming (2/3)

- Process each basic block in preorder of the dominator tree
-Rewrite each instruction (including the Phi nodes) in forward order
- For each use, replace the name with the latest name at the top of the stack
- For each def, generate a new name
- New name = original name + counter
- Increment the counter by 1
- Push the new name into the stack
-Propagate the new names to the Phi nodes of its successors
-Recursively process its children
-Pop names generated in this basic block from the stack



## Variable Renaming (2/3)

- Why preorder traversal
-If a variable has two definitions in different paths
- A Phi node would be inserted
- The two names for the definitions

-If a variable is defined only in the dominators
- The top of the stack is the name of the latest definition


## Variable Renaming: Example



```
var: a 
stk: a0 b0 c0 d0 i0
```


## Variable Renaming: Example



```
var: a 
stk: a0 b0 c0 d0 i0
```


## Variable Renaming: Example



```
var: a b c d i
ctr: 3 2 3 2 2
stk: a0 b0 c0 d0 i0
    a1 b1 c1 d1 il
    a2 c2
```


## Variable Renaming: Example



```
var: a b c d i
ctr: 3 3 4 4 3 2
stk: a0 b0 c0 d0 i0
    a1 b1 c1 d1 il
    a2 b2 c2 d2
        c3
```


## Variable Renaming: Example



```
var: a bllll
stk: a0 b0 c0 d0 i0
    a1 b1 c1 d1 il
    a2 c2
```

Pop names after BB2

## Variable Renaming: Example



```
var: a 
stk: a0 b0 c0 d0 i0
    a1 b1 c1 d1 il
    a2 c2 d3
    a3
```


## Variable Renaming: Example



| var: $a$ $b$ $c$ $d$ i <br> ctr: 4 3 4 5 2 <br> stk: a0 b0 c0 d0 <br> a1 i0    <br> a2  c1 d1 i1 <br> a3   d3  <br> a   d4  |  |  |  |
| :---: | :---: | :---: | :---: |
|  |  |  |  |
|  |  |  |  |
|  |  |  |  |
|  |  |  |  |

## Variable Renaming: Example



| var: $a$ | $b$ | $c$ | $d$ | $i$ |
| :--- | :--- | :--- | :--- | :--- |
| ctr: 4 | 3 | 5 | 5 | 2 |
| stk: | a0 | b0 | c0 | d0 |
| a1 | b1 | c1 | d1 | i1 |
| a2 |  | $c 2$ | d3 |  |
| a3 |  | $c 4$ |  |  |

## Variable Renaming: Example



```
var: a b c d i
ctr:4 4 6 6
stk: a0 b0 c0 d0 i0
    a1 b1 c1 d1 il
    a2 b3 c2 d3
    a3 c5 d5
```


## Variable Renaming: Example



```
var: a b c d i
ctr: 5 5 7 7 7 3
stk: a0 b0 c0 d0 i0
a1 b1 c1 d1 il
a2 b4 c2 d6 i2
    a4 c6
```


## Thanks \& all the best!

